

# An algebraic proof of the Frobenius condition for cubical sets

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Octoberfest 2019  
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# QMS from a premodel

## Definition

A *premodel* on a topos  $\mathcal{E}$  consists of  $(\Phi, \mathbb{I}, V)$  where:

- ▶  $\Phi$  is a representable class of monos  $\Phi \hookrightarrow \Omega^{\ast}$  ...
- ▶  $\mathbb{I}$  is an interval  $1 \rightrightarrows \mathbb{I}$  ...
- ▶  $\dot{V} \rightarrow V$  is a universe of *small families* ...

At CMU I sketched:

## Construction

*From a premodel  $(\Phi, \mathbb{I}, V)$  one can construct a QMS on  $\mathcal{E}$ .*

Today I will show that the resulting QMS is *right proper*.  
This only uses  $\Phi$  and  $\mathbb{I}$ .

## 1. The cofibration awfs $(\mathcal{C}, \text{TFib})$

The monos  $C \rightarrowtail Z$  classified by  $\Phi \hookrightarrow \Omega$  are the *cofibrations*  $\mathcal{C}$ .

$$\begin{array}{ccccc} C & \xrightarrow{\quad} & 1 & & \\ \downarrow & \nearrow & \downarrow & & \downarrow \\ & 1 & & & \Omega \\ \downarrow & \nearrow & \downarrow & & \downarrow \\ Z & \xrightarrow{\quad} & \Omega & & \end{array}$$

A commutative square diagram illustrating the cofibration awfs. The top row consists of objects  $C$ ,  $1$ , and  $\Omega$ . The bottom row consists of objects  $Z$  and  $\Omega$ . There are two vertical arrows: one from  $C$  to  $Z$ , and another from  $1$  to  $\Omega$ . There are two diagonal arrows: one from  $C$  to  $\Omega$  (the composition of the vertical and horizontal maps), and another from  $Z$  to  $1$  (the composition of the vertical and horizontal maps). The bottom-left arrow is labeled  $\Phi$ , and the bottom-right arrow is labeled  $\Omega$ . The middle vertical arrow is labeled  $1$ .

These are closed under pullbacks.

## 1. The cofibration awfs $(\mathcal{C}, \text{TFib})$

The generic cofibration  $1 \rightarrowtail \Phi$  determines a polynomial endofunctor,

$$X^+ := \sum_{\varphi:\Phi} X^\varphi.$$

This is a (fibered) monad,

$$+ : \mathcal{E} \longrightarrow \mathcal{E}$$

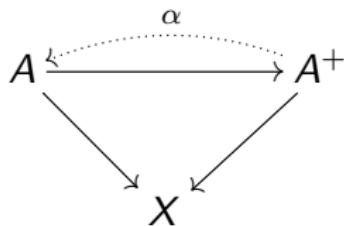
because of ...

## 1. The cofibration awfs $(\mathcal{C}, \text{TFib})$

In each slice  $\mathcal{E}/X$ , the algebras  $(A, \alpha)$  for the underlying pointed endofunctor,

$$+_X : \mathcal{E}/X \longrightarrow \mathcal{E}/X$$

are the *trivial fibrations*.



They form the right class of the *cofibration awfs*  $(\mathcal{C}, \text{TFib})$ .

## 1. The cofibration awfs $(\mathcal{C}, \text{TFib})$

The algebra structures on a trivial fibration correspond uniquely to *uniform right lifting structures* against the cofibrations:

$$\begin{array}{ccc} C & \longrightarrow & A \\ \downarrow & \nearrow \lrcorner & \downarrow \\ Z & \longrightarrow & X \end{array}$$

That is, a choice of fillers, ...

## 1. The cofibration awfs $(\mathcal{C}, \text{TFib})$

The algebra structures on a trivial fibration correspond uniquely to *uniform right lifting structures* against the cofibrations:

$$\begin{array}{ccccc} C' & \longrightarrow & C & \longrightarrow & A \\ \downarrow & \lrcorner & \downarrow & \lrcorner & \downarrow \\ Z' & \xrightarrow{\quad} & Z & \xrightarrow{\quad} & X \end{array}$$

That is, a choice of fillers, that are *coherent*.

## 1. The cofibration awfs $(\mathcal{C}, \text{TFib})$

The trivial fibrations are closed under all pullbacks, because they are “on the right”. But because their left class  $\mathcal{C}$  is also closed under all pullbacks, we have:

### Proposition

*The trivial fibrations are closed under all pushforwards.*

This follows by a standard “adjoint lemma”:

### Lemma

*Given a wfs  $(\mathcal{L}, \mathcal{R})$  and a base change  $f^* \dashv f_* : \mathcal{E}/X \longrightarrow \mathcal{E}/Y$ ,*

$f^*$  preserves  $\mathcal{L}$    iff    $f_*$  preserves  $\mathcal{R}$ .

## 2. The fibration awfs ( $\mathrm{TCof}, \mathcal{F}$ )

For any map  $u : A \rightarrow B$  in  $\mathcal{E}$ , the *Leibniz adjunction*

$$(-) \otimes u \dashv u \Rightarrow (-)$$

relates the pushout-product with  $u$  and the pullback-hom with  $u$ .

The functors  $(-) \otimes u \dashv (u \Rightarrow -) : \mathcal{E}^{\mathbb{D}} \longrightarrow \mathcal{E}^{\mathbb{D}}$  also satisfy

$$(c \otimes u) \boxtimes f \Leftrightarrow c \boxtimes (u \Rightarrow f)$$

with respect to the diagonal filling relation  $c \boxtimes f$ .

$$\begin{array}{ccc} \bullet & \xrightarrow{\quad} & \bullet \\ c \otimes u \downarrow & \nearrow \text{↗} & \downarrow f \\ \bullet & \xrightarrow{\quad} & \bullet \end{array} \qquad \begin{array}{ccc} \bullet & \xrightarrow{\quad} & \bullet \\ c \downarrow & \nearrow \text{↗} & \downarrow u \Rightarrow f \\ \bullet & \xrightarrow{\quad} & \bullet \end{array}$$

This holds also for uniform filling *structures*.

## 2. The fibration awfs ( $\text{TCof}, \mathcal{F}$ )

We then define the fibrations in terms of the trivial fibrations,

$$f \in \mathcal{F} \quad \text{iff} \quad \delta \Rightarrow f \in \text{TFib}$$

using the pullback-hom  $\delta \Rightarrow f$  with the *generic point*  $\delta : 1 \rightarrow \mathbb{I}$  in the slice category  $\mathcal{E}/\mathbb{I}$ .

### Definition

A map  $f : Y \rightarrow X$  is a *fibration* if  $\delta \Rightarrow f$  is a trivial fibration in  $\mathcal{E}/\mathbb{I}$ . Equivalently, by adjointness,  $f \in \mathcal{F}$  iff  $c \otimes \delta \dashv f$ ,

$$\begin{array}{ccc} Z +_C (C \times \mathbb{I}) & \xrightarrow{\quad} & Y \\ c \otimes \delta \downarrow & \nearrow & \downarrow f \\ Z \times \mathbb{I} & \xrightarrow{\quad} & X \end{array}$$

for all cofibrations  $c : C \rightarrow Z$ .

## 2. The fibration awfs ( $\mathrm{TCof}, \mathcal{F}$ )

A *fibration structure* on  $f : Y \rightarrow X$  thus corresponds to a *right lifting structure* against all  $c \otimes \delta$ ,

$$\begin{array}{ccc} \cdot & \xrightarrow{\quad} & Y \\ c \otimes \delta \downarrow & \nearrow & \downarrow f \\ Z \times \mathbb{I} & \xrightarrow{\quad} & X \end{array}$$

that is *uniform* with respect to all base changes  $u : Z' \rightarrow Z$ ,

$$\begin{array}{ccccc} \cdot & \xrightarrow{\quad} & \cdot & \xrightarrow{\quad} & Y \\ c' \otimes \delta \downarrow & \nearrow & \downarrow c \otimes \delta & \nearrow & \downarrow \\ Z' \times \mathbb{I} & \xrightarrow{\quad u \times \mathbb{I} \quad} & Z \times \mathbb{I} & \xrightarrow{\quad} & X \end{array}$$

## 2. The fibration awfs $(\text{TCof}, \mathcal{F})$

### Proposition

*There is an awfs  $(\text{TCof}, \mathcal{F})$  with these “uniform fibrations” as  $\mathcal{F}$ .*

Like the trivial fibrations, the fibrations are closed under pullbacks because they are “on the right”. But unlike the trivial fibrations, they are *not* closed under all pushforwards. However, we do have:

### Proposition

*The fibrations are closed under pushforward along fibrations.*

This is the main result to be shown below.

### 3. Frobenius

#### Proposition

*The fibrations are closed under pushforward along fibrations.*

Note that the “adjoint lemma” then implies:

#### Corollary (Frobenius)

*The trivial cofibrations are closed under pullback along fibrations.*

But since  $\mathcal{W} = \text{TFib} \circ \text{TCof}$ , and  $\text{TFib}$  is closed under pullbacks:

#### Proposition

$\mathcal{W}$  is closed under pullback along  $\mathcal{F}$ , i.e. the QMS is **right proper**.



### 3. Frobenius

Again, Frobenius says that trivial cofibrations pull back along fibrations,

$$\begin{array}{ccc} C' & \xrightarrow{\sim} & A \\ \downarrow & \lrcorner & \downarrow \\ C & \xrightarrow{\sim} & X. \end{array}$$

This is used for the interpretation of Id-types.

By the adjoint lemma, it is equivalent that fibrations push forward along fibrations,

$$\begin{array}{ccc} B & \twoheadrightarrow & A \\ & & \downarrow \\ \Pi_A B & \twoheadrightarrow & X. \end{array}$$

This is used for the interpretation of  $\Pi$ -types.

### 3. Frobenius

Recall that  $f : A \rightarrow X$  is a fibration iff  $\delta \Rightarrow f$  is a trivial fibration:

$$\begin{array}{ccccc} A^{\mathbb{I}} & & A^{\delta} & & \\ \downarrow \delta \Rightarrow f & & \searrow & & \\ & X^{\mathbb{I}} \times_X A & \xrightarrow{\quad} & A & \\ f^{\mathbb{I}} \swarrow & & \downarrow & & \downarrow f \\ & X^{\mathbb{I}} & \xrightarrow{\quad} & X^{\delta} & \end{array}$$

We indicate this briefly as follows:

$$\begin{array}{ccccc} A^{\mathbb{I}} & \xrightarrow{\delta \Rightarrow A} & A_{\epsilon} & \longrightarrow & A \\ & \searrow & \downarrow & & \downarrow \\ & & X^{\mathbb{I}} & \xrightarrow{\epsilon} & X \end{array}$$

### 3. Frobenius

#### Proposition

*Fibrations push forward along fibrations.*

#### Proof.

Consider fibrations  $B \rightarrow A \rightarrow X$ .

Thus we have:

$$\begin{array}{ccccc} B^{\mathbb{I}} & \xrightarrow{\delta \rightarrow B} & B_{\epsilon}^* & \longrightarrow & B \\ \searrow & & \downarrow & & \downarrow \\ & & A^{\mathbb{I}} & \xrightarrow{\delta \rightarrow A} & A_{\epsilon} \\ & & \searrow & & \downarrow \\ & & X^{\mathbb{I}} & \xrightarrow{\epsilon} & X \end{array}$$

### 3. Frobenius

Taking the pushforward of the right column yields  $\Pi_A B \rightarrow X$ .

$$\begin{array}{ccccc}
B^{\mathbb{I}} & \xrightarrow{\delta \Rightarrow B} & B_{\epsilon}^* & \longrightarrow & B \\
\searrow & & \downarrow & & \downarrow \\
& & A^{\mathbb{I}} & \xrightarrow{\delta \Rightarrow A} & A_{\epsilon} \\
& & \searrow & & \downarrow \\
& & X^{\mathbb{I}} & \xrightarrow{\epsilon} & X \\
\uparrow & & \nearrow & & \uparrow \\
(\Pi_A B)^{\mathbb{I}} & \xrightarrow{\delta \Rightarrow \Pi_A B} & (\Pi_A B)_{\epsilon} & \longrightarrow & \Pi_A B
\end{array}$$

We want to show that  $\delta \Rightarrow \Pi_A B$  is a trivial fibration.

### 3. Frobenius

The pushforward of  $\delta \Rightarrow B$  along  $A^{\mathbb{I}} \rightarrow X^{\mathbb{I}}$  is a trivial fibration over  $X^{\mathbb{I}}$ , since these are closed under all pushforwards.

$$\begin{array}{ccccc}
B^{\mathbb{I}} & \xrightarrow{\delta \Rightarrow B} & B_{\epsilon}^* & \longrightarrow & B \\
\searrow & & \downarrow & & \downarrow \\
& & A^{\mathbb{I}} & \xrightarrow{\delta \Rightarrow A} & A_{\epsilon} \\
& & \searrow & & \downarrow \\
& & X^{\mathbb{I}} & \xrightarrow{\epsilon} & X \\
\swarrow & & \uparrow f' & & \uparrow \\
(\Pi_A B)^{\mathbb{I}} & \xrightarrow{\delta \Rightarrow \Pi_A B} & (\Pi_A B)_{\epsilon} & \longrightarrow & \Pi_A B \\
& & \uparrow & & \uparrow \\
\Pi_{A^{\mathbb{I}}} B^{\mathbb{I}} & \xrightarrow{\Pi_{A^{\mathbb{I}}}. \delta \Rightarrow B} & \Pi_{A^{\mathbb{I}}} B_{\epsilon}^* & &
\end{array}$$

### 3. Frobenius

One then shows that there is a retraction of  $\Pi_{A^{\mathbb{I}}}. \delta \Rightarrow B$  onto  $\delta \Rightarrow \Pi_A B$  over  $X^{\mathbb{I}}$ .

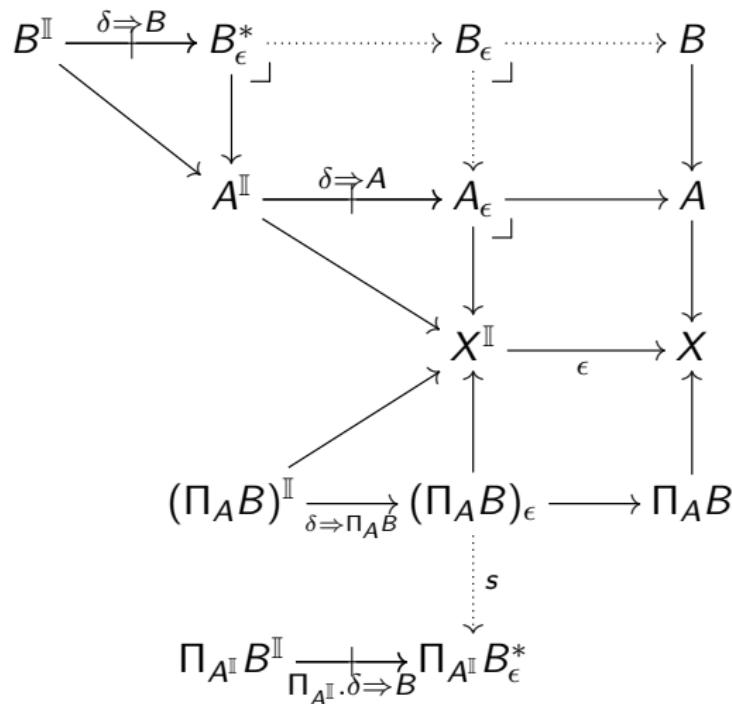
$$\begin{array}{ccccc}
B^{\mathbb{I}} & \xrightarrow{\delta \Rightarrow B} & B_{\epsilon}^* & \longrightarrow & B \\
\searrow & & \downarrow & & \downarrow \\
& & A^{\mathbb{I}} & \xrightarrow{\delta \Rightarrow A} & A_{\epsilon} \\
& & \searrow & & \downarrow \\
& & X^{\mathbb{I}} & \xrightarrow{\epsilon} & X \\
& & \uparrow & & \uparrow \\
(\Pi_A B)^{\mathbb{I}} & \xrightarrow{\delta \Rightarrow \Pi_A B} & (\Pi_A B)_{\epsilon} & \longrightarrow & \Pi_A B \\
\text{↗} \quad \text{↘} & & \text{↗} \quad \text{↘} & & \\
\text{---} \quad \text{---} & & \text{---} \quad \text{---} & & \\
\Pi_{A^{\mathbb{I}}} B^{\mathbb{I}} & \xrightarrow{\Pi_{A^{\mathbb{I}}}. \delta \Rightarrow B} & \Pi_{A^{\mathbb{I}}} B_{\epsilon}^* & &
\end{array}$$

□

### 3. Frobenius

For example, to get  $s : (\Pi_A B)_\epsilon \rightarrow \Pi_{A^{\mathbb{I}}} B_\epsilon^*$  first interpolate  $B_\epsilon$  so

$$(\Pi_A B)_\epsilon \cong \Pi_{A_\epsilon} B_\epsilon$$



### 3. Frobenius

Then use  $\Pi_{A^{\mathbb{I}}} B_{\epsilon}^* \cong \Pi_{A_{\epsilon}}(B_{\epsilon}^*)_*$

$$\begin{array}{ccccccc}
 & & & (B_{\epsilon}^*)_* & & & \\
 & & & \uparrow \eta & & & \\
 B^{\mathbb{I}} & \xrightarrow{\delta \Rightarrow B} & B_{\epsilon}^* & \dashrightarrow & B_{\epsilon} & \dashrightarrow & B \\
 & \searrow & \downarrow & & \downarrow & & \downarrow \\
 & & A^{\mathbb{I}} & \xrightarrow{\delta \Rightarrow A} & A_{\epsilon} & \longrightarrow & A \\
 & & \searrow & & \downarrow & & \downarrow \\
 & & X^{\mathbb{I}} & \xrightarrow{\epsilon} & X & & \\
 & \nearrow & \uparrow & & \uparrow & & \\
 (\Pi_A B)^{\mathbb{I}} & \xrightarrow{\delta \Rightarrow \Pi_A B} & (\Pi_A B)_{\epsilon} & \longrightarrow & \Pi_A B & & \\
 & & & & \downarrow s := \Pi_{A_{\epsilon}} \eta & & \\
 & & \Pi_{A^{\mathbb{I}}} B^{\mathbb{I}} & \xrightarrow[\Pi_{A^{\mathbb{I}}}. \delta \Rightarrow B]{} & \Pi_{A^{\mathbb{I}}} B_{\epsilon}^* & & 
 \end{array}$$



## 4. References

This “algebraic” proof of Frobenius is derived from a type theoretic one due to Thierry Coquand:

- ▶ Cohen, Coquand, Huber, Mörtberg:  
Cubical Type Theory: A constructive interpretation of the univalence axiom, TYPES 2015.

Also see:

- ▶ Awodey: A Quillen model structure on cartesian cubical sets,  
[github.com/awodey/math/qms](https://github.com/awodey/math/qms) (2019)